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STRUCTURAL COMPLETENESS AND THE DISJUNCTION PROPERTY OF INTERMEDIATE LOGICS

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Let At be the set of all propositional variables p_0, p_1, p_2, \dots , let S be the set of well-formed formulas built by means of variables p_0, p_1, p_2, \dots and connectives: $\rightarrow, *, +, \sim$. By intermediate logic we mean a set $H \subseteq S$ such that $INT \subseteq H \Rightarrow Cn(H) \subseteq L_2$, where INT is the intuitionistic propositional logic, L_2 – two-valued logic and $Cn(H)$ is the smallest set containing H and closed under the substitution rule and the detachment rule.

Kreisel and Putnam [1] have proved that the logic KP obtained by adding to INT the axiom $(p \rightarrow q + r) \rightarrow (\sim p \rightarrow q) + (\sim p \rightarrow r)$ has the following disjunction property:

$$\Phi + \Psi \in KP \Leftrightarrow \Phi \in KP \vee \Psi \in KP.$$

In this paper it is shown that there exist a structural complete intermediate logics with the disjunction property, which was previously conjectured by H. Friedman.

The intermediate logic H is structurally complete ($H \in SCpl$) if and only if for every $\Phi, \Psi \in S$ the following condition holds:

$$\forall e: At \rightarrow S [h^e(\Phi) \in H \Rightarrow h^e(\Psi) \in H] \Leftrightarrow (\Phi \rightarrow \Psi) \in H,$$

where h^e is the extension of the function $e: At \rightarrow S$ to the endomorphism of the algebra $\langle S, \rightarrow, *, +, \sim \rangle$, (the notion of structural completeness is introduced by W. A. Pogorzelski in [2]).

Let
 $\widetilde{At} = \{\sim p \mid p \in At\};$
 S = the least set containing At and closed with respect to: $\rightarrow, *, +, \sim;$
 $H = \{\Phi \in S \mid \forall_{e: At \rightarrow S} h^e(\Phi) \in H\}$, for every intermediate logics with the disjunction property.

We have

THEOREM. *For every intermediate logic H :*

- i) $H \in SCpl \Rightarrow KP \subseteq H$
- ii) $KP \subseteq H \Rightarrow \widetilde{H} \in SCpl$
- iii) $(KP \subseteq H \text{ and } H \in GP) \Rightarrow \widetilde{H} \in GP.$

COROLLARY. $\widetilde{KP} \in SCpl \cap GP.$

References

- [1] G. Kreisel, H. Putnam, *Eine Unableitbarkeitsbeweismethode für den intuitionistischen Aussagenkalkül*, **Arch. Math. Logik Grundlagenforsch.** 3 (1957), pp. 74–78.
- [2] W. A. Pogorzelski, *Structural completeness of the propositional calculus*, **Bull. Acad. Polon. Sci.**, Sér. Sci. Math. Astronom. Phys., 19 (1971), pp. 349–351.

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