

Wiesław Dziobiak

SEMANTICS OF KRIPKE'S STYLE FOR SOME MODAL SYSTEMS*

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We use the well-known and set theoretical notation. The logical connectives will be represented by \sim, \rightarrow, L, M denote negation, material implication, necessity and possibility, respectively. Formulas are represented by the capitals A, B, C, \dots ; the symbols k, n, i, \dots denote the natural numbers $1, 2, 3, \dots$.

Here is list of formulas important for our purpose.

- (1) $LA \rightarrow A$
- (2ⁿ) $L^n A \rightarrow L^{n+1} A$
- (3ⁿ) $M^n L^n A \rightarrow L^n A$
- (4ⁿ) $L^n A \rightarrow M^n L^{n+1} A$
- (5ⁿ) $L^n A \rightarrow M^n A$
- (6ⁿ) $L^n A \rightarrow M^n L^n A$

Let Cn denote a consequences operation determined by the modal system K of Kripke, the detachment and Gödel's rule. The present paper gives of Kripke's style for the following modal systems:

$$\begin{aligned} S5_n &= Cn((1), (2^n), (3^n)) \\ S4_n &= Cn((1), (2^n)) \\ T_n^k &= Cn((1), (2^k), (4^n)), \quad (k \geq n) \end{aligned}$$

*As abstract this article is not to be reviewed.

$$\begin{aligned} T_n^* &= Cn((1), (4^n)) \\ C_n^* &= Cn((4^n), (5^n), (6^n)) \end{aligned}$$

The Kripke's style semantics for these systems will be fixed in order to use them for an explanation of some syntactic relation holding between these systems (see [1], [2], [3], [4]). Note that $S4_1$ and $S5_1$ are the well known modal systems $S4$ and $S5$ of Lewis; $S4_n$ ($n \geq 2$) constituting the family of Sobociński's modal systems. T_n and T_n^k have been considered in [1], [2], [4].

By a model structure we understand any pair $\langle W, R \rangle$, where W is a non-empty set and $R \subset W \times W$. Assuming $A \subset W$ and $R \subset W \times W$ we define

$$\begin{aligned} R|_A &= \{(w, w') \in W \times W : (w, w') \in A \wedge ((w, w') \in R)\} \\ R^{-1} &= \{(w, w') \in W \times W : (w', w) \in R\} \\ R^n &= \{(w, w') \in W \times W : \exists_{\{w_i\}_{i \leq n-1}} ((w, w_1) \in R) \wedge ((w_1, w_2) \in R) \dots ((w_{n-1}, w) \in R)\} \end{aligned}$$

The set $\{(w, w) \in W \times W : w \in W\}$ is denoted by Δ . For any model structure $\langle W, R \rangle$ and any $w \in W$ by a *tree* with top w we understand the set T_w having properties:

- (i) $w \in T_w$
- (ii) if $(w', w'') \in R$ and $w' \in T_w$, then $w'' \in T_w$.

THEOREM 1. *$S5_n$ is determined by the class of all model structures $\langle W, R \rangle$ provided R satisfies:*

- (i) $\Delta \subset R$
- (ii) $R^{n+1} \subset R^n$
- (iii) $R^n \subset (R^n)^{-1}$

The similar theorem for $S4_n$ looks as follows.

THEOREM 2. *$S4_n$ is determined by the class of all model structures $\langle W, R \rangle$ provided R satisfies:*

- (i) $\Delta \subset R$
- (ii) $R^{n+1} \subset R^n$
- (iii) $\forall_w \exists_{w'} ((w, w') \in R^n) \wedge (\forall_{w'' \in T_{w'}} (w, w'') \in R^n) \wedge (R^n|_{T_{w'}} \subset (R^n|_{T_{w'}})^{-1})$

Note that Theorem 2 supplies some other semantics for the systems from Sobociński's family $S4_n$ and the modal system $S4$ of Lewis.

THEOREM 3. T_n^k ($k \geq n$) is determined by the class of all model structures $\langle W, R \rangle$ provided R satisfies:

- (i) $\Delta \subset R$
- (ii) $R^{k+1} \subset R^k$
- (iii) $\forall_w \exists_{w'} ((w, w') \in R^n) \wedge (\forall_{w'' \in T_w} (w, w'') \in R^n) \wedge (R_{|T_{w'}}^{n+1} \subset R_{|T_{w'}}^n) \wedge \wedge (R_{|T_{w'}}^n \subset (R_{|T_{w'}}^n)^{-1})$

THEOREM 4. T_n^* is determined by the class of all model structures $\langle W, R \rangle$ provided R satisfies:

- (i) $\Delta \subset R$
- (ii) $\forall_w \exists_{w'} (((w, w') \in R^n) \wedge (\forall_{w'' \in T_w} (w, w'') \in R^n) \wedge (R_{|T_{w'}}^{n+1} \subset R_{|T_{w'}}^n) \wedge \wedge (R_{|T_{w'}}^n \subset (R_{|T_{w'}}^n)^{-1}))$

THEOREM 5. C_n^* is determined by the class of all model structures $\langle W, R \rangle$ provided R satisfies:

- i) $\Delta \subset R$
- ii) $\forall_w \exists_{w'} (((w, w') \in R^n) \wedge (\forall_{w'' \in T_w} (w, w'') \in R^n) \wedge \wedge (R_{|T_{w'}}^{n+1} \subset R_{|T_{w'}}^n) \wedge \wedge (R_{|T_{w'}}^n \subset (R_{|T_{w'}}^n)^{-1}))$

For any modal system S , denote by $M^n - S$ the set of all formulas which, when preceded by M (n -times), become thesis of S . Using results contained in [1], [2], [3], [4] we obtain that between the systems T_n^* , T_n^k , $S4_n$, $S5_n$ the following relation occurs:

$$(1) \quad M^n - T_n^* = M^n - T_n^k = M^n - S4_n = M^n - S5_n.$$

Notice that the given semantics of Kripke's style for T_n^* , T_n^k , $S4_n$, $S5_n$ confirm the validity of the condition (1).

Resting on the semantics of Kripke's style for C_n^* (Th. 5) one can prove the following

$$\text{THEOREM 6. } M^n - C_n^* = M^n - S5_n.$$

Note that from Theorems 3 and 5 we have $C_n^* \subsetneq T_n$. The following theorem states that C_n^* are the least modal systems in the class of all normal

modal systems containing the Kripke's modal system K and having the property: $M^n - C_n^* = M^n - S5_n$.

THEOREM 7. For any normal modal system $S \supseteq X$:
If $M^n - S = M^n - S5_n$, then $C_n^* \subseteq S$.

References

- [1] J. J. Błaszczyk, W. Dziobiak, Modal systems related to $S4_n$ of Sobociński, this **Bulletin**, vol. 4, no. 3 (1975), pp. 103–108.
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Chair of Logic
Nicholas Copernicus University
Toruń