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ON FINITELY BASED CONSEQUENCE OPERATIONS

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S. L. Bloom [2] posed the following two questions (the first of them can also be found in R. Wójcicki [8]):

- (1) whether the consequence operation determined by a finite matrix is always finitely based,
- (2) whether the meet of two finitely based consequence operations is always finitely based.

In this paper both of these questions are answered negatively.

§0. Let $\mathbb{T} = \langle T, d_0, \dots, d_n \rangle$ be a free algebra of terms of some finite type free generated by an infinite set of variables $V = \{x_0, x_1, \dots\}$. We shall use Greek lower case α, β, \dots for terms and Greek upper case Φ, Ψ, \dots for sets of terms. The symbol $V(\Phi)$ denotes the set of variables occurring in terms of Φ . If X is a set then the symbols $P(X), P_f(X)$ denote the set of subsets of X and the set of finite subsets of X respectively.

A mapping $C : P(T) \rightarrow P(T)$ is a consequence operation in T iff for every $\Phi, \Psi \subseteq T$, $\Phi \subseteq CC(\Phi) \subseteq C(\Phi \cup \Psi)$. A consequence operation C is algebraic (or finite) iff for every $\Phi \subseteq T$, $C(\Phi) = \bigcup CP_f(\Phi)$, a consequence operation C is structural iff for every $\Phi \subseteq T$, $\varepsilon \in \text{Hom}(\mathbb{T}, \mathbb{T})$, $\varepsilon C(\Phi) \subseteq C(\varepsilon\Phi)$. If C is structural and algebraic then it is called standard.

A rule is a subset $\varrho \subseteq P_f(T) \times T$, a rule ϱ is structural iff it is closed under substitutions i.e. for every $\varepsilon \in \text{Hom}(\mathbb{T}, \mathbb{T})$, if $\langle \Phi, \alpha \rangle \in \varrho$ then $\langle \varepsilon\Phi, \varepsilon(\alpha) \rangle \in \varrho$. A sequent is a pair $\langle \Phi, \alpha \rangle \in P_f(T) \times T$. Every sequent $\langle \Phi, \alpha \rangle$ determines a structural rule $\varrho(\Phi, \alpha) = \{ \langle \varepsilon\Phi, \varepsilon(\alpha) \rangle : \varepsilon \in \text{Hom}(\mathbb{T}, \mathbb{T}) \}$.

A rule ϱ is called sequential if $\varrho = \varrho(\Phi, \alpha)$ for some sequent $\langle \Phi, \alpha \rangle$. Every set of rules R determines a consequence operation C_R such that for every $\Phi \subseteq T$, $C_R(\Phi)$ is the smallest subset of T containing Φ and closed under rules of R . A consequence operation C is finitely based iff $C = C_R$ for some finite set R of sequential rules. It should be noted that the concept of finitely based consequence operation is to be credited to R. Suszko [5] who introduced it under the name of finitely formalizable consequence operation. In the sequel we shall use the following well-known results:

THEOREM 0.1. (Łoś, Suszko [4]) *A consequence operation C is standard iff $C = C_R$ for some countable set R of sequential rules.*

THEOREM 0.2. (Bloom [1]) *A standard consequence operation C is not finitely based iff $C = \sup\{C_n\}_{n < \omega}$ for some strictly increasing sequence $C_0 < C_1 < \dots$ of standard consequence operations.*

A matrix is a pair $\mathcal{M} = \langle \mathbb{M}, D \rangle$ where \mathbb{M} is an algebra similar to \mathbb{T} and D is a subset of the carrier of \mathbb{M} . Every matrix \mathcal{M} determines a consequence operation $C_{\mathcal{M}}$ such that $\alpha \in C_{\mathcal{M}}(\Phi)$ iff for every valuation $\vartheta \in \text{Hom}(\mathbb{T}, \mathbb{M})$, if $\vartheta\Phi \subseteq D$ then $\vartheta(\alpha) \in D$. Let us note the following:

THEOREM 0.3. (Łoś, Suszko [4]) *If \mathcal{M} is a finite matrix then the consequence operation $C_{\mathcal{M}}$ is standard.*

Following J. Zygmunt [9] we say that a matrix \mathcal{M} is proper iff for every $\vartheta \in \text{Hom}(\mathbb{T}, \mathbb{M})$, $\vartheta T \not\subseteq D$. Observe that if \mathcal{M} is proper then $C_{\mathcal{M}}(\Phi) \neq T$ iff $\vartheta\Phi \subseteq D$ for some $\vartheta \in \text{Hom}(\mathbb{T}, \mathbb{M})$. This observation yields the following fact being an immediate consequence of a theorem of [9].

THEOREM 0.4. (comp. Zygmunt [9]) *Let $\mathcal{M}_1, \mathcal{M}_2$ be proper matrices, then for every $\Phi \subseteq T$ the following condition holds:*

$$C_{\mathcal{M}_1 \times \mathcal{M}_2}(\Phi) = \begin{cases} T & \text{if } C_{\mathcal{M}_1}(\Phi) = T \text{ or } C_{\mathcal{M}_2}(\Phi) = T, \\ C_{\mathcal{M}_1}(\Phi) \cap C_{\mathcal{M}_2}(\Phi) & \text{otherwise.} \end{cases}$$

§1. In this section we shall give an example of finite matrix whose consequence operation is not finitely based. From now on $\mathbb{T} = \langle T, \cdot \rangle$ is assumed to be the free algebra of terms of the type $\langle 2 \rangle$ free generated by an infinite set of variables $V = \{x_0, x_1, \dots\}$. In order to simplify the notations we adopt the convention of associating to the left and ignoring

the symbol of binary operation, thus for example we shall write $\alpha\beta\gamma\alpha(\alpha\alpha)$ instead of $((\alpha \cdot \beta) \cdot \gamma) \cdot \alpha) \cdot (\alpha \cdot \alpha)$.

Let $\mathbb{A} = \langle \{0, 1, 2\}, \cdot \rangle$ be an algebra of the type $\langle 2 \rangle$ whose binary operation \cdot is such that $0 \cdot 0 = 2 \cdot 2 = 2$, $1 \cdot 1 = 1$ and $a \cdot b = 0$ otherwise. Note that the set $\{0, 2\}$ is a subuniverse of \mathbb{A} and the subalgebra $\mathbb{B} = \langle \{0, 2\}, \cdot \rangle$ of \mathbb{A} is the two-element equivalential algebra (see [3]). Define a matrix $\mathcal{A} = \langle \mathbb{A}, \{0\} \rangle$ and a matrix $\mathcal{Z} = \langle \mathbb{B}, \{0\} \rangle$. It is easy to see that the matrices \mathcal{A}, \mathcal{Z} are proper and \mathcal{Z} is a submatrix of \mathcal{A} . We shall prove the following:

PROPOSITION 1.1. *The consequence operation $C_{\mathcal{A} \times \mathcal{Z}}$ is not finitely based.*

In preparation for the proof we need to state several lemmas. First let us observe that the consequence operation $C_{\mathcal{A}}$ is standard by virtue of Theorem 0.3 and therefore applying Theorem 0.1 we can choose a countable set R of sequential rules such that $C_{\mathcal{A}} = C_R$. For every $n = 1, 2, \dots$ we define a set of sequential rules $R_n = \{\rho(\Phi, \alpha) : \Phi \in P_f(T), |V(\Phi)| \leq n, \alpha \in T = C_{\mathcal{Z}}(\Phi)\}$ and the corresponding consequence operation $C_n = C_{R \cup R_n}$. By Theorem 0.1 it follows that each the consequence operation C_n is standard and moreover we have the following:

LEMMA 1.1. $C_{\mathcal{A}} \leq C_n < C_{n+1}$.

PROOF. Since $R \subseteq R \cup R_n \subseteq R \cup R_{n+1}$ then $C_{\mathcal{A}} \leq C_n \leq C_{n+1}$ and it remains to show that $C_n \neq C_{n+1}$. For every $n = 1, 2, \dots$ we define a term $\gamma_n = (x_0 x_0)(x_1 x_1) \dots (x_n x_n)$. It is easy to see that $C_{\mathcal{Z}}(\gamma_n) = T$ which means that $\xi(\{\gamma_n\}, x_{n+1}) \in R_{n+1}$ and consequently $C_{n+1}(\gamma_n) = T$. On the other hand $C_{\mathcal{A}}(\gamma_n) \neq T$ (consider a valuation $\vartheta_0 \in \text{Hom}(\mathbb{T}, \mathbb{A})$ such that $\vartheta_0(x_0) = 1$ and $\vartheta_0(x_i) \neq 1$ whenever $i \neq 0$) and therefore it suffices to show that $C_n(\gamma_n) \subseteq C_{\mathcal{A}}(\gamma_n)$. We shall prove that if $|V(\Phi)| \leq n$ and $C_{\mathcal{Z}}(\Phi)$ then for every $\varepsilon \in \text{Hom}(\mathbb{T}, \mathbb{T})$, $\varepsilon\Phi \notin C_{\mathcal{A}}(\gamma_n)$ which means that $C_{\mathcal{A}}(\gamma_n)$ is closed under rules of R_n . Suppose that $|V(\Phi)| \leq n$, $C_{\mathcal{Z}}(\Phi) = T$ and $\varepsilon \in \text{Hom}(\mathbb{T}, \mathbb{T})$. Since $|V(\gamma_n)| = n + 1$ then $|V(\Phi)| < |V(\gamma_n)|$ and therefore there must exist a variable $x_c \in V(\gamma_n)$ such that for every variable $x_j \in V(\Phi)$, $\{x_c\} \neq V(\varepsilon(x_j))$. This gives that for every $x_j \in V(\Phi)$, $x_c \notin V(\varepsilon(x_j))$ or $\{x_c\} \subsetneq V(\varepsilon(x_j))$. Let $\vartheta_c \in \text{Hom}(\mathbb{T}, \mathbb{A})$ be such that $\vartheta_c(x_c) = 1$ and $\vartheta_c(x_i) \neq 1$ whenever $i \neq c$. Then for every $x_j \in V(\Phi)$ $\vartheta_c(\varepsilon(x_j)) \in \{0, 2\}$ because every term $\varepsilon(x_j)$ must contain a variable whose value under ϑ_c is distinct from 1. Pick a valuation $\bar{\vartheta}_c \in \text{Hom}(\mathbb{T}, \mathbb{B})$ such that $\bar{\vartheta}_c(x_j) = \vartheta_c(\varepsilon(x_j))$ for every $x_j \in V(\Phi)$. Then $\bar{\vartheta}_c(\alpha) = \vartheta_c(\varepsilon(\alpha))$

whenever $V(\alpha) \subseteq V(\Phi)$ which yields that $\vartheta_c \varepsilon \Phi = \bar{\vartheta}_c \Phi \not\subseteq \{0\}$ because $C_{\mathcal{Z}}(\Phi) = T$ and the matrix \mathcal{Z} is proper. The observation that $\vartheta_c(\gamma_n) = 0$ completes the proof of lemma. Q.E.D.

Note that from Lemma 1.1 follows the existence of a strictly increasing sequence of standard consequence operations beginning with a consequence operation determined by a finite matrix. Such a result (disproving a conjecture of R. Wójcicki [7]) was first obtained by M. Tokarz. The present proof of this result is a modification of the basic idea of M. Tokarz [6].

Let $C_\omega = \sup\{C_n\}_{n < \omega}$, then $C_\omega = C_{R \cup \bigcup\{R_n\}_{n < \omega}}$ and from Lemma 1.1 and Theorem 0.2 it follows that C_ω is not finitely based. Thus Proposition 1.1 can be obtained as a direct consequence of the following:

LEMMA 1.2. $C_{\mathcal{A} \times \mathcal{Z}} = C_\omega$.

§2. In this section it will be shown that the meet of two finitely based consequence operations need not to have any finite base.

Let us define two sequential rules: $r = \{\langle \{\alpha\}, \alpha\beta \rangle : \alpha, \beta \in T\}$ and $l = \{\langle \{\alpha\}, \beta\alpha \rangle : \alpha, \beta \in T\}$ and the corresponding two consequence operations Cr and Cl in T determined by the rules r and l respectively. We shall prove the following:

PROPOSITION 2.1. *The consequence operation $\inf\{Cr, Cl\}$ is not finitely based.*

First, for every $n = 0, 1, \dots$ we define two mappings $Cr^n, Cl^n : P(T) \rightarrow P(T)$ putting for every $\Phi \subseteq T$, $\alpha \in T$: $\alpha \in Cr^n(\Phi)$ iff for some $m \leq n$ there exists a sequence of terms $\Delta_0, \dots, \Delta_m$ such that $\Delta_m = \alpha$ and for every $i \leq m$, if $\Delta_i \in \Phi$ then $\Delta_i = \Delta_j \beta$ for some $j < i$, $\beta \in T$; $\alpha \in Cl^n(\Phi)$ iff for some $m \leq n$ there exists a sequence of terms $\Delta_0, \dots, \Delta_m$ such that $\Delta_m = \alpha$ and for every $i \leq m$, if $\Delta_i \notin \Phi$ then $\Delta_i = \beta \Delta_j$ for some $j < i$, $\beta \in T$.

The proof of the following simple lemma is left to the reader:

LEMMA 2.1.

- (i) $Cr^n(\Phi) \subseteq Cr^{n+1}(\Phi) \subseteq Cr(\Phi)$, $Cl^n(\Phi) \subseteq Cl^{n+1}(\Phi) \subseteq Cl(\Phi)$;
- (ii) $Cr(\Phi) = \bigcup\{Cr^n(\Phi)\}_{n < \omega}$, $Cl(\Phi) = \bigcup\{Cl^n(\Phi)\}_{n < \omega}$;
- (iii) $Cr(\Phi) = \bigcup\{Cr(\alpha) : \alpha \in \Phi\}$, $Cl(\Phi) = \bigcup\{Cl(\alpha) : \alpha \in \Phi\}$.

For every $n = 1, 2, \dots$ we define a matrix $\mathcal{N}_n = \langle \mathbb{N}_n, \{0\} \rangle$ where $\mathbb{N}_n =$

$\langle \{0, \dots, n\}, \cdot \rangle$ is an algebra of the type $\langle 2 \rangle$ whose binary operation \cdot is such that $(n-1) \cdot 0 = n$, $a \cdot (a+1) = a+1$ whenever $a+1 < n$ and $a \cdot b = 0$ otherwise.

LEMMA 2.2. *If $n = 1, 2, \dots$, $\vartheta \in \text{Hom}(\mathbb{T}, \mathbb{N}_n)$ and α is a term such that $\vartheta(\alpha) = 0$ then $\vartheta Cl(\alpha) \subseteq \{0, n\}$ and $\vartheta Cr^{n-1}(\alpha) \subseteq \{0, \dots, n-1\}$.*

Now for every $n = 0, 1, \dots$ we define a set of sequential rules $Q_n = \{\rho(\{\alpha, \beta\}, \gamma) : \gamma \in Cr^n(\alpha) \cap Cl^n(\beta)\}$ and the corresponding consequence operation $Cq_n = C_{Q_n}$.

LEMMA 2.3. $Cq_n < Cq_{n+1}$.

Let $Cq_\omega = \sup\{Cq_n\}_{n < \omega}$, then $Cq_\omega = C_{\bigcup\{Q_n\}_{n < \omega}}$ and from Lemma 2.3 and Theorem 0.2 it follows that Cq_ω is not finitely based. Thus Proposition 2.1 follows immediately from the following:

LEMMA 2.4. $Cq_\omega = \inf\{Cr, Cl\}$.

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