

Marek Tokarz

DEDUCTION THEOREMS FOR RM AND ITS EXTENSIONS

This is an abstract of a lecture read at the Seminar of the Logic Department, Wrocław University, February 1977.

Consider the following Sugihara algebra $\underline{S} = (Q, \rightarrow, \vee, \wedge, \sim)$, where Q is the set of all the rational numbers (including 0) and the operations $\rightarrow, \vee, \wedge, \sim$ are defined as follows:

$$\begin{aligned}\sim x &= -x \\ x \vee y &= \max(x, y) \\ x \wedge y &= \min(x, y) \\ x \rightarrow y &= \begin{cases} -x \vee y & \text{if } x \leq y, \\ -x \wedge y & \text{otherwise.} \end{cases}\end{aligned}$$

Let $\underline{S}_a, \underline{S}_2, \underline{S}_3, \underline{S}_4, \dots$ be subalgebras of \underline{S} generated by the sets $\{0\}, \{-1, 1\}, \{-1, 0, 1\}, \{-2, -1, 1, 2\}, \dots$ respectively. Put $D = \{x \in Q : x \geq 0\}$. Then the pair (\underline{S}, D) is a Sugihara matrix with D as the set of designated elements. We shall also use the matrices $(\underline{S}_1, D_1), (\underline{S}_2, D_2), \dots$ where $D_i = D \cap |\underline{S}_i|$. The symbol \underline{S}_i will denote both the i -element algebra as well as the suitable i -element matrix.

Let $\underline{L} = (L, \rightarrow, \vee, \wedge, \sim)$ be a propositional language with p_1, p_2, \dots as propositional variables. Let (\underline{A}, B) be any matrix under consideration. By $C_{(\underline{A}, B)}$ we shall understand the matrix consequence in \underline{L} in the sense of Łoś and Suszko [2], i.e. for every $X \cup \{\alpha\} \subseteq \underline{L}$

$$\alpha \in C_{(\underline{A}, B)}(X) \text{ iff } \forall h : \underline{L} \rightarrow^{hom} \underline{A} (hX \subseteq B \Rightarrow h\alpha \in B).$$

To make formulas short we shall write C_1 instead of $C_{\underline{S}_1}$. Let us define additional consequences in \underline{L} , namely C_{RM} and C_i^0 by putting:

- $C_{RM}(X) =$ the least set of formulas including $X \cup RM$ and closed under modus ponens and adjunction.
 $C_i^0(X) =$ the least set of formulas including $X \cup C_i(\emptyset)$ and closed under modus ponens and adjunction.

The most important results on RM and Sugihara matrices are included in Dunn's [1], it is presupposed here that paper is familiar to the reader.

By Ackermann's rule we mean the one given by the scheme $\alpha, \sim \alpha \vee \beta / \beta$. Moreover, we use new connectives \supset and \succ defined in the following way:

$$\begin{aligned} \alpha \supset \beta &=_{df} [\sim (\alpha \rightarrow \sim \beta) \vee (\alpha \rightarrow \beta)] \wedge (\sim \alpha \vee \beta), \\ \alpha \succ \beta &=_{df} \beta \vee (\alpha \rightarrow \beta). \end{aligned}$$

The following theorems hold:

1. (Dunn [1])
 - a. $C_{RM} = C_S$.
 - b. The only systems including RM and closed under substitution, modus ponens and adjunction are those of the form $C_i(\emptyset)$.
 - c. $C_1(\emptyset) \supsetneq C_2(\emptyset) \supsetneq C_3(\emptyset) \supsetneq \dots$
 - d. $RM = \bigcap \{C_i(\emptyset) : i \in \omega\}$.
2. $\beta \in C_{2i}(X \cup \{\alpha\})$ iff $\sim \alpha \vee \beta \in C_{2i}(X)$.
3. $\beta \in C_3(X \cup \{\alpha\})$ iff $\sim (\alpha \rightarrow \sim \beta) \vee (\alpha \rightarrow \beta) \in C_3(X)$.
4. $\beta \in C_{2i+1}(X \cup \{\alpha\})$ iff $\alpha \supset \beta \in C_{2i+1}(X)$.
5. $C_{2i+1} = C_{2i+1}^0$.
6. If $i \geq 2$, then $C_{2i} > C_{2i}^0$.
7. Ackermann's rule is not deducible in C_{2i}^0 , all $i \geq 2$.
8. $\beta \in C_{2i}^0(X \cup \{\alpha\})$ iff $\alpha \succ \beta \in C_{2i}^0(X)$.
9. $\beta \in C_{RM}(X \cup \{\alpha\})$ iff $\alpha \supset \beta \in C_{RM}(X)$.

References

- [1] J. M. Dunn, *Algebraic completeness results for R-mingle and its extensions*, **The Journal of Symbolic Logic** 35 (1970), pp. 1–13.

[2] J. Łoś and R. Suszko, *Remarks on sentential logics*, **Indagationes Mathematicae** 20 (1958), pp. 177–183.

*The Section of Logic
Institute of Philosophy and Sociology
Polish Academy of Sciences*