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DEGREES OF MAXIMALITY OF ŁUKASIEWICZ-LIKE SENTENTIAL CALCULI

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Łukasiewicz-like n -valued sentential calculi

The algebra of formulas

$$\underline{L} = (L, \rightarrow, \vee, \wedge, \neg) \quad (1)$$

will be called the *language of Łukasiewicz-like sentential calculi* ($\rightarrow, \vee, \wedge, \neg$ denote the implication, disjunction, conjunction and negation connectives, respectively). Now consider the algebras

$$\underline{A}_n = (A_n, \rightarrow, \vee, \wedge, \neg) \quad (2)$$

where n is an arbitrary natural number, $n \geq 2$, similar to \underline{L} and defined as follows:

- (i) $A_n = \{0, \frac{1}{n-1}, \dots, \frac{n-2}{n-1}, 1\}$
- (ii) $x \rightarrow y = \min(1, 1 - x + y)$; $x \vee y = \max(x, y)$
 $x \wedge y = \min(x, y)$; $\neg x = 1 - x$
for $x, y \in A_n$

Given $n \geq 2$, let $I \subsetneq A_n$ be any subset of A_n such that $1 \in I$ and $0 \notin I$. Then the matrix

$$M_n^I = (\underline{A}_n, I) \quad (3)$$

will be called the *n-valued Lukasiewicz-like matrix of the type I*. The symbol C_n^I will be used to denote the consequence operation determined by the matrix M_n^I (in \underline{L}).

The pair

$$\mathbf{L}_n^I = (\underline{L}, C_n^I) \quad (4)$$

will be called the *n-valued Lukasiewicz-like sentential calculus of the type I*. Clearly $M_n^{\{1\}} = M_n$, where M_n is the well known *n*-valued Lukasiewicz matrix and $\mathbf{L}_n^{\{1\}} = \mathbf{L}_n$ appears to be the *n*-valued Lukasiewicz calculus.

Given $n \geq 2$, the class of all *n*-valued Lukasiewicz-like sentential calculi is a superclass of Rosser Turquette's [2] class of generalizations of the *n*-valued Lukasiewicz calculus. Recall that for every matrix M_n^I in [2] $I = \{x \in A_n : x \geq i_{n-1}\}$, where i_{n-1} is some given element of A_n not equal to 0.

2. \mathbf{L}_n^I -algebras

For the whole present section let us assume that $\mathbf{L}_n^I = (\underline{L}, C_n^I)$ is some fixed *n*-valued Lukasiewicz-like sentential calculus. For every $x, y \in A_n$ let us put

$$x \Rightarrow y = \begin{cases} 1 & \text{if } x \leq y \\ 0 & \text{otherwise.} \end{cases} \quad (5)$$

Clearly, \Rightarrow is definable in M_n^I , cf. [2]. The same symbol \Rightarrow will be used to denote an appropriate sentential connective in \underline{L} .

Now, let us consider matrices $M = (\underline{A}_M, I_M)$, where $\underline{A}_M = (A_M, \rightarrow, \vee, \wedge, \neg)$ are algebras similar to \underline{L} and $I_M \subseteq A_M$. Given such M , by Cn_M we shall denote the consequence operation determined by M (in \underline{L}). The symbol $\text{Matr}(C_n^I)$ will be used to denote the class of all C_n^I -matrices, i.e. the class of matrices M , such that $C_n^I \leq Cn_M$.

For every $M \in \text{Matr}(C_n^I)$ we put

$$a \approx_M b \text{ if and only if } a \Rightarrow b, b \Rightarrow a \in I_M. \quad (6)$$

and

$$V_M = \{a \in A_M : \text{there exists a formula } \alpha \in L \text{ such that} \quad (7)$$

$j_1(\alpha) \in C_n^I(\emptyset)$ and such that there exists a

homomorphism $h : \underline{L} \rightarrow \underline{A}_M$ for which $h\alpha = a\}$,

where j_1 is a sentential connective corresponding to Rosser and Turquette's j_1 ($j_1(x) = 1$ whenever $x = 1$ and $j_1(x) = 0$ otherwise), cf. [2].

LEMMA 1.

- (i) The relation \approx_M is a congruence of M
- (ii) V_{M/\approx_M} is a one element set
- (iii) Put $1_M = V_{M/\approx_M}$ (see (ii)). Then if $|a| \rightarrow |b| = 1_M$ and $|b| \rightarrow |a| = 1_M$, then $|a| = |b|$.

In the sequel we shall consider the class

$$\text{Matr}^R(C_n^I) = \{M/\approx_M : M \in \text{Matr}(C_n^I)\}. \quad (8)$$

3. The equational characterization of $\text{Matr}^R(C_n^I)$

Given finite $n \geq 2$, let $HSP(\underline{A}_n)$ denote the smallest equational class of algebras containing the algebra \underline{A}_n . Now let us denote by $HSP(M_n^I)$ a corresponding class of matrices defined in the following manner: For every $\underline{A}_M \in HSP(\underline{A}_n)$ we take the matrix $M = (\underline{A}_M, I_M)$, where I_M is determined by one of the following conditions:

- (i₁) If $\underline{A}_M = \underline{A}_M$, then $I_M = I$
- (i₂) If \underline{A}_M is a product of the indexed set of algebras $\{\underline{A}_{M_i} : i \in T\}$, then $I_M = \prod\{I_{M_i} : i \in T\}$
- (i₃) If \underline{A}_M is a subalgebra of some algebra \underline{A}_{M_0} , then $I_M = \underline{A}_M \cap I_{M_0}$
- (i₄) If $\underline{A}_M = h(\underline{A}_{M_0})$, where h is a homomorphism, then $I_M = h(I_{M_0})$.

In the sequel, the element of $\underline{A}_M \in HSP(\underline{A}_n)$ being a natural counterpart of $1 \in A_n$ will be denoted by 1_M (cf. Lemma 1).

Let us now assume that $I = \{m_1, \dots, m_k\}$, and consider Rosser and Turquette's functions j_{m_1}, \dots, j_{m_k} , i.e. such functions of \underline{A}_n , that

$$j_{m_i}(x) = \begin{cases} 1 & \text{if } x = m_i \\ 0 & \text{otherwise,} \end{cases} \quad (9)$$

cf. [2]. Because j_{m_i} 's are definable in \underline{A}_n , we can assume $j_{m_i}(\alpha)$, where $\alpha \in L$, to be formulas of \underline{L} .

In turn, let us put

$$\varphi_I(\alpha) = j_{m_1}(\alpha) \vee \dots \vee j_{m_k}(\alpha); \quad (10)$$

obviously $\varphi_I(\alpha) \in L$. Moreover, for any $x \in A_n$ we have

$$\varphi_I(x) = \begin{cases} 1 & \text{if } x \in I \\ 0 & \text{otherwise.} \end{cases} \quad (11)$$

From the representation theorem for $HSP(\underline{A}_n)$ and the properties of φ_I , we obtain

LEMMA 2. $M = (\underline{A}_n, I_M) \in HSP(M_n^I)$ if and only if 1^0 . $\underline{A}_M \in HSP(\underline{A}_n)$ and 2^0 . $I_M = \{a \in A_M; \varphi_I(a) = 1_M\}$.

In turn, using the results of [3] and Lemma 2, one can prove

LEMMA 3. $HSP(M_n^I) \subseteq Matr^R(C_n^I)$.

Now, we are going to construct Lindenbaum matrix for \mathbf{L}_n^I . In the language \underline{L} we introduce the following relation:

$$\alpha \approx \beta \text{ if and only if } \alpha \Rightarrow \beta, \beta \Rightarrow \alpha \in C_n^I(\emptyset). \quad (12)$$

One can easily verify that \approx is a congruence relation on L . The quotient matrix

$$\Lambda_n = (\underline{L}/\approx, Cn_n^I(\emptyset)/\approx) \quad (13)$$

will be called the *Lindenbaum matrix* for \mathbf{L}_n^I .

LEMMA 4. \underline{L}/\approx is free in the class $\{\underline{A}_M : (\underline{A}_M, I_M) \in Matr(C_n^I) \text{ for some } I_M \subseteq A_M\}$, the set $\{\|p\| : p \in V\}$ (where V is the set of all sentential

variables in \underline{L}) being the set of free generators in \underline{L}/\approx . If $p_1 \neq p_2$, then $\|p_1\| \neq \|p_2\|$.

Using the last lemma and Lemma 2, one can prove the following

LEMMA 5. $Matr^R(C_n^I) \subseteq HSP(M_n^I)$.

Combining the results of Lemma 3 and Lemma 5 we obtain the main theorem of the present section.

THEOREM 1. $Matr^R(C_n^I) = HSP(M_n^I)$.

4. Degrees of maximality of \mathbf{L}_n^I

Given a sentential calculus (\underline{L}, C) , \underline{L} being a sentential language C a structural consequence operation on \underline{L} , the *degree of maximality* of (\underline{L}, C) , $dm(\underline{L}, C)$, is equal to the cardinal number of all structural consequence operations C' such that $C \leq C'$ (cf. e.g. [4]). Let us now consider any Łukasiewicz-like calculus \mathbf{L}_n^I . By the argument in [4] we obtain

$$dm(\mathbf{L}_n^I) = \text{card}\{Cn_M : M \subseteq Matr(C_n^I)\}. \quad (14)$$

Let $M \in Matr(C_n^I)$ be an arbitrary matrix and let \approx_M be the congruence relation defined by (6). From the fact that \approx_M is a congruence of M (Lemma 1 (i)), it follows that

$$Cn_M = Cn/\approx_M \quad (15)$$

As it was defined in Section 2, the whole class of such M/\approx_M is $Matr^R(C_n^I)$. Moreover, according to the equational characterization given in Section 3 we have $Matr^R(C_n^I) = HSP(M_n^I)$. This implies that the sequence operation determined by the matrix belonging to the class $HSP(M_n^I)$. Consequently,

$$dm(\mathbf{L}_n^I) = \text{card}\{Cn_M : M \subseteq HSP(M_n^I)\}. \quad (16)$$

Using Lemma 9 in [4] and a correspondence between $HSP(M_n^I)$ and $HSP(M_n)$ (cf. Lemma 2), one can prove the following:

LEMMA 6. *For each $M \in HSP(M_n^I)$ there are pairwise different submatrices N_1, \dots, N_s of M_n^I such that $Cn_M = Cn_{N_1 \times \dots \times N_s}$.*

In turn, the following theorem is an easy corollary to (16) and Lemma 6.

THEOREM 2. *The degree of maximality of any (finite) Łukasiewicz-like sentential calculus L_n^I is finite.*

REMARK. A more close inspection of the classes $HSP(M_n^I)$ and $HSP(M_n)$ leads to the conclusion that for any given $n \geq 2$, the degrees of maximality of the calculi L_n and L_n^I are equal. In particular for the calculi L_n, L_n^I for which $(n - 1)$ is prime, the degree of maximality equals 4 (cf. [1]).

References

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