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## A CHARACTERIZATION OF $Matr(C)$

This is an abstract of the paper “Reduced products of logical matrices” submitted to *Studia Logica*.

Propositional languages are defined as in the note [2]. If  $\underline{L}$  and  $\underline{L}'$  are two propositional languages and  $\underline{L}$  is a subalgebra of  $\underline{L}'$  then  $\underline{L}$  will be called a sublanguage of  $\underline{L}'$  and  $\underline{L}'$  an extension of  $\underline{L}$ , symbolically  $\underline{L} \subseteq \underline{L}'$ . Notice that a subalgebra of  $\underline{L}'$  need not be a sublanguage of  $\underline{L}'$ .

Let  $C$  be a consequence operation on a language  $\underline{L}$ . By the cardinality of  $C$ ,  $card(C)$ , we shall mean the least cardinal  $n$  such that for all  $X \subseteq L$

$$C(X) = \bigcup \{C(Y) \mid Y \subseteq X \text{ \& } \bar{Y} < n\}$$

$C$  is finite iff  $card(C) < \aleph_0$ .

Given two consequence operations  $C_1, C_2$  on  $\underline{L}$  we write  $C_1 \leq C_2$  iff  $C_1(X) \subseteq C_2(X)$ , all  $X \subseteq L$ .

If  $\mathcal{M} = (\underline{A}, D)$  is a matrix for  $\underline{L}$  then  $Cn_{\mathcal{M}, \underline{L}}$  is the consequence operation on  $\underline{L}$  induced by  $\mathcal{M}$ . If  $\mathbb{K}$  is a class of matrices for  $\underline{L}$  then

$$Cn_{\mathbb{K}, \underline{L}} =_{df} \inf_{\mathcal{M} \in \mathbb{K}} Cn_{\mathcal{M}, \underline{L}}$$

where the infimum is taken with respect to  $\leq$ .

If  $C$  is consequence operation on  $\underline{L}$  then we define  $Matr(C)$  to be the class of all matrices  $\mathcal{M}$  for  $\underline{L}$  such that  $C \leq Cn_{\mathcal{M}, \underline{L}}$ .

A consequence operation  $C$  on  $\underline{L}$  is *structural* provided that  $\varepsilon(C(X)) \subseteq C(\varepsilon(X))$ , all  $X \subseteq L$ , all endomorphisms  $\varepsilon$  of  $\underline{L}$ .

By an  $\mathfrak{m}$ -filter over a set  $I$  we shall mean a filter  $\mathcal{F}$  closed with respect to intersections of fewer than  $\mathfrak{m}$  members of  $\mathcal{F}$ . Every improper filter  $\mathcal{F}$  is an  $\mathfrak{m}$ -filter.

Let  $\mathbb{K}$  be a class of matrices for  $\underline{L}$ . We use the following terminology:

$S(\mathbb{K})$  – the class of all isomorphic images of submatrices of members of  $\mathbb{K}$ ,

$P(\mathbb{K})$  – the class of all isomorphic images of direct products of arbitrary systems (possibly empty) of members of  $\mathbb{K}$ ,

$P_{\mathfrak{m}-r}(\mathbb{K})$  – the class of all isomorphic images of  $\succ$ -reduced products (i.e. direct products modulo  $\mathfrak{m}$ -filters) of arbitrary systems of members of  $\mathbb{K}$ ,

$H_s(\mathbb{K})$  – the class of all strong homomorphic images of members of  $\mathbb{K}$ ,

$\overleftarrow{H}_s(\mathbb{K})$  – the class of all strong homomorphic counter – images of members of  $\mathbb{K}$ . Hence  $\mathcal{M} \in \overleftarrow{H}_s(\mathbb{K})$  iff there exists a strong homomorphism from  $\mathcal{M}$  onto a matrix in  $\mathbb{K}$ . Clearly  $\mathbb{K} \subseteq H_s(\mathbb{K})$  and  $\mathbb{K} \subseteq \overleftarrow{H}_s(\mathbb{K})$ .

If  $\mathfrak{m} = \aleph_0$  then we write

$P_r(\mathbb{K})$  instead of  $P_{\aleph_0-r}(\mathbb{K})$ .

**THEOREM 1.** *Let  $C$  be a structural consequence operation on a propositional language  $\underline{L}$ . Let  $\{\mathcal{M}_i\}_{i \in I}$  be a family of matrices for  $\underline{L}$  such that  $C \leq Cn_{\mathcal{M}_i, \underline{L}}$ , all  $i \in I$ . Let  $\mathfrak{m}$  be an infinite cardinal  $\geq \text{card}(C)$  and let  $\mathfrak{m}$  be an  $\mathfrak{m}$ -filter over  $I$ .*

*Form the  $\mathfrak{m}$ -reduced product  $\mathcal{M} = \prod_{i \in I} \mathcal{M}_i$ . Then  $C \leq Cn_{\mathcal{M}, \underline{L}}$ .*

**THEOREM 2.** *Let  $\mathbb{K}$  be a class of matrices for a propositional language  $\underline{L}$ . Let  $\mathfrak{m}$  be a regular infinite cardinal such that  $\text{card}(Cn_{\mathbb{K}, \underline{L}}) \leq \mathfrak{m} \leq \bar{\bar{L}}^+$ . Then*

$$\text{Matr}(Cn_{\mathbb{K}, \underline{L}}) = \overleftarrow{H}_s H_s S P_{\mathfrak{m}-r}(\mathbb{K}).$$

For a given structural  $C$  in  $\underline{L}$  let

$Th(C) =_{df} \{X \mid X \subseteq L \text{ \& } C(X) = X\}$  and

$\mathbb{L}_C =_{df} \{\underline{L}_X \mid \underline{L}_X = (\underline{L}, X) \text{ \& } X \in Th(C)\}$ .

$\mathbb{L}_C$  is called the *Lindenbaum bundle* of  $C$ . As known  $C = Cn_{\mathbb{L}_C, \underline{L}}$  [1]. Let  $\mathbb{L}_C^* = \{\underline{L}_{X/\theta_X} \mid X \in Th(C)\}$ , where  $\theta_X$  is the greatest congruence is the matrix  $\underline{L}_X$  (see [1]).

**THEOREM 3.** *Let  $C$  be a structural consequence operation on a propositional language  $\underline{L}$ . Let  $\mathfrak{m}$  be an infinite regular cardinal such that  $\text{card}(C) \leq \mathfrak{m} \leq \bar{\bar{L}}^+$ . Then*

$$\text{Matr}(C) = \overleftarrow{H}_s S P_{\mathfrak{m}-r}(\mathbb{L}_C^*).$$

**THEOREM 4.** *Let  $\mathbb{K}$  be a class of matrices for a propositional language  $\underline{L}$ .*

The following are equivalent:

- (a)  $\text{Matr}(Cn_{\mathbb{K}, \underline{L}}) = \overleftarrow{H}_s H_s SP(\mathbb{K})$
- (b) For every extension  $\underline{L}'$  of  $\underline{L}$

$$\text{Matr}(Cn_{\mathbb{K}, \underline{L}}) \subseteq \text{Matr}(Cn_{\mathbb{K}, \underline{L}'}).$$

**THEOREM 5.** Let  $\mathbb{K}$  be a finite family of finite matrices for a propositional language  $\underline{L}$ . Then

$$\text{Matr}(Cn_{\mathbb{K}, \underline{L}}) = \overleftarrow{H}_s H_s SP(\mathbb{K}).$$

For a given class  $\mathbb{K}$  of matrices for  $\underline{L}$  ( $\underline{L}$ -fixed) define

$$\mathcal{M} \sim \mathcal{N} \text{ iff } Cn_{\mathcal{M}, \underline{L}} = Cn_{\mathcal{N}, \underline{L}}$$

$$\mathcal{M} \precsim \mathcal{N} \text{ iff } Cn_{\mathcal{M}, \underline{L}} \leq Cn_{\mathcal{N}, \underline{L}}$$

( $\mathcal{M}, \mathcal{N} \in \mathbb{K}$ ). Then  $\mathbb{K}/\sim$  is a set and  $\mathbb{K}/\sim$  is partially ordered by  $\precsim$ .

**THEOREM 6.** Let  $\mathbb{K}$  be a finite class of matrices for a propositional language  $\underline{L}$ . Then the partially ordered sets  $\langle \text{Matr}(Cn_{\mathbb{K}, \underline{L}})/\sim, \precsim \rangle$  and  $\langle SP(\mathbb{K})/\sim, \precsim \rangle$  are isomorphic.

## References

- [1] R. Wójcicki, *Matrix approach in methodology of sentential calculi*, **Studia Logica**, Vol. XXXII 1973.
- [2] J. Czelakowski, 'Large' matrices which induce finite consequence operations, this **Bulletin**, pp. 79–82.

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