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OMEGA-CONSISTENCY AND THE DIAMOND (Abstract)

G is the result of adjoining the schema $\Box(\Box A \rightarrow A) \rightarrow \Box A$ to K ; the axioms of G^* are the theorems of G and the instances of the schema $\Box A \rightarrow A$ and the sole rule of G^* is modus ponens. A sentence is ω -provable if it is provable in P(eano) A(rithmetic) by one application of the ω -rule; equivalently, if its negation is ω -inconsistent in PA. Let $\omega\text{-Bew}(x)$ be the natural formalization of the notion of ω -provability. For any modal sentence A and function ϕ mapping sentence letters to sentences of PA, inductively define $A^{\omega\phi}$ by: $p^{\omega\phi} = \phi(p)$ (p a sentence letter); $\perp^{\omega\phi} = \perp$; $(A \rightarrow B)^{\omega\phi} = (A^{\omega\phi} \rightarrow B^{\omega\phi})$; and $(\Box A)^{\omega\phi} = \omega\text{-Bew}(\ulcorner A^{\omega\phi} \urcorner)$ ($\ulcorner S \urcorner$ is the numeral for the Gödel number of the sentence S). Then, applying techniques of Solovay (**Israel Journal of Mathematics** 25, pp. 287–304), we prove that for every modal sentence A , $\vdash_G A$ iff for all ϕ , $\vdash_{PA} A^{\omega\phi}$; and for every modal sentence A , $\vdash_{G^*} A$ iff for all ϕ , $A^{\omega\phi}$ is true. (To appear in **Studia Logica**).

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