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ON THE VERIFIERS OF DISJUNCTION

In a letter David Makinson has pointed out to me a clear-cut counterexample to the formula T3: $S(\alpha \vee \beta) = \text{Min}(S(\alpha) \cup S(\beta))$, propounded in [1] for objective of disjunction. As he says: let p, q be independent E -propositions and put $\alpha = p \wedge q$, $\beta = p \wedge \sim q$. Since p is an E -proposition, $S(p) = \{u\}$ for some u in SE . Then $S(\alpha \vee \beta) = S((p \wedge q) \vee (p \wedge \sim q)) = S(p) = \{u\}$; but $u \notin V(\alpha) \cup V(\beta) \supset S(\alpha) \cup S(\beta) \supset \text{Min}(S(\alpha) \cup S(\beta))$.

The formula T3 was entailed (cf. [2]) by the axiom A5: $V(\alpha \vee \beta) = V(\alpha) \cup V(\beta)$. Consequently, the essential half of A5 has to be dropped too, though by A2 we are still left with the other one:

$$/D1/ \quad V(\alpha) \cup V(\beta) \subset V(\alpha \vee \beta).$$

However, in view of the other axioms, the following equivalence holds with regard to the verifiers of disjunction:

$$/D2/ \quad s \in V(\alpha \vee \beta) \equiv \bigwedge_{t \geq s} \bigvee_{u \geq t} u \in V(\alpha) \cup V(\beta).$$

To see this observe that

$$/N/ \quad s \in V(\sim \alpha) \equiv \sim \bigvee_t (s \leq t \wedge t \in V(\alpha))$$

is a theorem. Indeed, implication \leftarrow of N is equivalent to A9. On the other hand, if $s \in V(\sim \alpha)$ and $s \leq t$, then $t \in V(\sim \alpha)$ by A3. Hence $\sim t \in V(\alpha)$ by A6, cancelling double negation. Starting now with $V(\alpha \vee \beta) = V(\sim (\sim \alpha \wedge \sim \beta))$, D2 may be derived from N by A4, exactly as the theorem 2.4/iii

is in [3], if only we read Bell's " p forces σ " as " $p \in V(\sigma)$ ", and reverse everywhere his " \leq ".

Let D_1, D_2, \dots, D_n be the SE -sets corresponding to the logical dimensions of a W -language L . (Their number is assumed to be finite!) Thus $D_1 \cup D_2 \cup \dots \cup D_n = \text{Min}(SE)$, and the product $PL = D_1 \cdot D_2 \cdot \dots \cdot D_n = \text{Max}(SE)$ is then the *logical space* of L , i.e. it is the totality of *logical points* or possible worlds of L . (Cf. [4].) We have here the lemma: $V(\alpha) \cdot PL = V(\alpha) \cap PL$. (The inclusion \supset is obvious. Conversely, if $u = s$; $w \wedge s \in V(\alpha) \wedge w \in PL$, then $u \in V(\alpha)$ by A3, and $u \in PL$ by w being maximal). This leads to the theorem:

$$/D3/ \quad V(\alpha \vee \beta) \cdot PL \subset V(\alpha) \cup V(\beta),$$

for suppose $w \in V(\alpha \vee \beta)$ and $w \in PL$. Then $w \in V(\alpha) \cup V(\beta)$ by D2, which in view of the lemma yields D3.

We adopt now a new axiom:

$$/A10/ \quad A \cdot PL \subset V(\alpha) \rightarrow A \subset V(\alpha)$$

for any $A \subset SE$. (We preserve the numbering of [1], though A5 is wanting; moreover, A7 and A8 are redundant, since A8 follows from A7, and A7 is provable with the number of dimensions being finite). From this axiom we get another formula for $V(\alpha \vee \beta)$:

$$/D4/ \quad V(\alpha \vee \beta) = \bigcup \{A \subset SE; A \cdot PL \subset V(\alpha) \cup V(\beta)\}.$$

Indeed, suppose $s \in V(\alpha \vee \beta)$. Then by D3 we have $s \in A \wedge A \cdot PL \subset V(\alpha) \cup V(\beta)$ for some A , i.e. $s \in \bigcup \{\dots\}$. Conversely, if $s \in \bigcup \{\dots\}$, then by D1: $s \in A \wedge A \cdot PL \subseteq V(\alpha \vee \beta)$, for some $A \subseteq SE$. Hence by A10 we get $s \in V(\alpha \vee \beta)$.

But what about the objective $S(\alpha \vee \beta)$ as a function of $V(\alpha)$ and $V(\beta)$? This will be discussed in a wider context, together with the objectives of tautology and negation.

References

- [1] B. Wolniewicz, *Objectives of Propositions*, this **Bulletin**, vol. 7, no. 3 (1978), pp. 143–147.
- [2] B. Wolniewicz, *Some Formal Properties of Objectives*, this **Bulletin**, vol. 8, no. 1 (1979), pp. 16–20.
- [3] J. L. Bell, **Boolean-Valued Models and Independence Proofs**, Oxford 1977.
- [4] R. Suszko, *Ontology in the Tractatus of L. Wittgenstein*, **Notre Dame Journal of Formal Logic**, vol. 9, no. 1 (1968), p. 15.

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