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## SOME EXTENSIONS OF THE BROUWERIAN LOGIC

The Brouwerian logic ( $B$ ) is obtained by adding  $Lp \supset p$  and  $p \supset LMp$  to the minimal normal logic,  $K$ .  $B + Lp \supset LLp$  yields  $S5$ , and  $B + p \supset Lp$  the ‘trivial’ modal logic (Triv). All the above *wff* are of the form  $Xp \supset Yp$ , where  $X$  and  $Y$  are affirmative modalities (possibly empty). This paper discusses various other logics obtained by adding such *wff* to  $B$ .

By ‘modality’ we mean ‘affirmative modality’ throughout. We use  $X, Y, Z$  to range over modalities, and  $\alpha, \beta$  over *wff*. By  $X'$  we mean the result of replacing  $L$  by  $M$  and  $M$  by  $L$  everywhere in  $X$ .  $L^n$  and  $M^n$  designate sequences of  $n$   $L$ ’s and  $n$   $M$ ’s respectively. If  $X$  consists of precisely  $k$  operators it is said to be of *degree*  $k$ .

The following rules and theorem schemata (where  $\alpha$  and  $\beta$  are any *wff* and  $X$  and  $Y$  are any modalities) hold in all normal extensions of  $B$ :

- $R1.$   $\vdash \alpha \rightarrow \vdash X\alpha$
- $R2.$   $\vdash \alpha \supset \beta \rightarrow \vdash X\alpha \supset X\beta$
- $R3.$   $\vdash Xp \supset Yp \rightarrow \vdash Y'p \supset X'p$
- $T1.$   $\vdash X\alpha \supset XL^nM^n\alpha$  ( $n \geq 0$ )
- $T2.$   $\vdash XM^nL^n\alpha \supset X\alpha$  ( $n \geq 0$ )
- $T3.$   $\vdash L^{n+m}\alpha \supset L^n\alpha$  ( $n, m \geq 0$ )
- $T4.$   $\vdash M^n\alpha \supset M^{n+m}\alpha$  ( $n, m \geq 0$ )
- $T5.$  If  $X$  is of degree  $k$ ,  $\vdash X\alpha \supset M^k\alpha$
- $T6.$  If  $n \geq m$ ,  $\vdash (L^n\alpha \supset M^m\beta) \supset M^n(\alpha \supset \beta)$
- $T7.$   $\vdash M^n(\alpha \supset \beta) \equiv (L^n\alpha \supset M^n\beta)$  ( $n \geq 0$ )

A modality containing no sequence of the form  $M^nL^nM^n$  or  $L^nM^nL^n$  ( $n \geq 1$ ) will be said to be *in reduced form* (or simply, *reduced*). Since  $\vdash_B M^nL^nM^n\alpha = M^n\alpha$  and  $\vdash_B L^nM^n\alpha = L^n\alpha$  ( $n \geq 1$ ), we can without loss of generality assume that all modalities are in reduced form.

The best-known extensions of  $B$  are the  $T_n^+$  logics. For each  $n \geq 0$ ,  $T_n^+$  is defined as  $B + T^n p \supset L^{n+1} p$ . (Thus  $T_1^+$  is  $S5$  and  $T_0^+$  is  $Triv$ ). We note that in each  $T_n^+$  the following theorem schema holds:

$$T8. \quad \vdash X\alpha \supset Y M^n \alpha$$

A  $wff \alpha$  will be said to be a  $B^*$  formula iff either  $\vdash_B \alpha$  or  $B + \alpha = T_n^+$  for some  $n$ .

LEMMA 1. For every  $n \geq 0$ , if  $\alpha$  is of the form  $p \supset XL^{n+1} M^n p$  then  $B + \alpha = T_n^+$ .

PROOF. By  $T8$ ,  $T_n^+$  contains  $B + \alpha$ . For the converse, assume  $B$  and  $p \supset XL^{n+1} M^n p$ . Then by  $[L^n p/p]$  we have

$$\begin{aligned} (1) \quad & L^n p \supset XL^{n+1} M^n L^n p \\ (2) \quad & L^n p \supset XL^{n+1} p \quad ((1), T2) \end{aligned}$$

Let  $k$  be the degree of  $X$ . Then

$$\begin{aligned} (3) \quad & L^n p \supset M^k L^{n+1} p \quad ((2), T5) \\ (4) \quad & L^k M^{n+1} p \supset M^n p \quad ((3), R3) \end{aligned}$$

Since, by  $R1$ ,  $\vdash M^{n+1} \alpha \rightarrow \vdash L^k M^{n+1} \alpha$ , (4) yields the rule:  $\vdash M^{n+1} \alpha \rightarrow \vdash M^n \alpha$ ; and hence by repetition the rule:

$$\begin{aligned} (5) \quad & \vdash M^n \alpha \rightarrow \vdash M^n \alpha \quad \text{for every } m \geq n \\ (6) \quad & L^{2n+k} p \supset M^k L^{2n+k+1} p \quad ((3)[L^{n+k} p/p]) \\ (7) \quad & M^{2n+k} (p \supset L^{2n+k+1} p) \quad ((6), T6) \\ (8) \quad & M^n (p \supset L^{2n+k+1} p) \quad ((7), (5)) \\ (9) \quad & L^n p \supset M^n L^{2n+k+1} p \quad ((8), T7) \\ (10) \quad & \supset L^{n+k+1} p \quad ((9), T2) \\ (11) \quad & \supset L^{n+1} p \quad ((10), T3) \end{aligned}$$

COROLLARY. If  $Z$  ends with  $L$ , then  $B + p \supset Zp = Triv$ ; and if  $Z$  ends with  $LIM$ , then  $B + p \supset Zp = S5$ .

LEMMA 2. If  $Z$  is a reduced modality not of the form  $XL^{n+1} M^n$  then  $Z$  is of the form

$$L^{a_1} M^{a_2} \dots L^{a_{n-1}} M^{a_n}$$

where  $a_1, \dots, a_n$  are non-negative integers such that (i)  $a_{n-1} \leq a_n$  and (ii)  $a_1 < \dots < a_{n-1}$ .

PROOF. Clearly every modality is of the form mentioned, where  $a_1, \dots, a_n$  are *some* non-negative integers. The Lemma then follows from the easily verified facts (a) that if  $Z$  violates (i) then it is of the form  $XL^{n+1}M^n$ ; and (b) that if it violates (ii) but not (i) then it contains some sequence  $M^{a_j}L^{a_{j+1}}M^{a_{j+2}}$  or  $L^{a_j}M^{a_{j+1}}L^{a_{j+2}}$  such that  $a_j \geq a_{j+1} \leq a_{j+2}$ ; and hence is not in reduced form.

LEMMA 3. *If  $Z$  is a reduced modality not of the form  $XL^{n+1}M^n$  then  $\vdash_B p \supset Zp$ .*

The proof is straightforward from Lemma 2, T1 and T4.

Since, as noted above, every modality is equivalent in  $B$  to some reduced modality, Lemmas 1 and 3 immediately yield.

THEOREM 1. *Every wff of the form  $p \supset Zp$  is a  $B^*$  formula.*

Clearly, to show that  $\alpha$  is a  $B^*$  formula it is sufficient to show that, given  $B$ ,  $\alpha$  is interdeducible with some wff of the form  $p \supset Zp$ .

THEOREM 2. *Every wff of the form  $M^jL^kp \supset Yp$  ( $i, k \geq 0$ ) is a  $B^*$  formula.*

PROOF. Assume  $B$  and  $M^jL^kp \supset Yp$ . Then by  $[M^kp/p]$  and  $R2(X = L^j)$  we have  $L^jM^jL^kM^kp \supset L^jYM^kp$ . Whence by T1 we have  $p \supset L^jYM^kp$ , which is of the form  $p \supset Zp$ . Now assume  $B$  and  $p \supset L^jYM^kp$ . Then by  $[L^kp/p]$  and  $R2(X = M^j)$  we obtain  $M^jL^kp \supset M^jL^jYM^kL^kp$ . Whence by T2 we have  $M^jL^kp \supset Yp$ .

THEOREM 3. *Every wff of the form  $L^jM^kp \supset YM^lp$  ( $j \geq k \leq l$ ) is a  $B^*$  formula.*

PROOF. Assume  $B$  and (1)  $L^jM^kp \supset YM^lp$ . Then since  $j \geq k$ , we have

$$\begin{array}{ll} (2) & L^jM^jp \supset YM^{l+j-k}p \quad ((1)[M^{j-k}p/p]) \\ (3) & p \supset YM^{l+j-k}p \quad ((2), T1) \end{array}$$

(3) is of the form  $p \supset Zp$ , and hence by Theorem 1 is a  $B^*$  formula. Now assume  $B$  and (3). Then

$$(4) \quad L^jM^kp \supset YM^{l+j-k}L^jM^kp \quad ((3)[L^jM^kp/p])$$

Since  $k \leq l$ , we have

- $$\begin{array}{ll}
(5) & M^{l-k}M^jp \supset M^{l-k}p \quad (T2, R2 \ (X = M^{l-k})) \\
i.e. (6) & M^{l+j-k}L^jp \supset M^{l-k}p \\
(7) & YM^{l+j-k}L^jM^kp \supset YM^lp \quad ((6)[M^kp/p], R2) \\
(1) & L^jM^kp \supset YM^lp \quad ((4), (6), Syll)
\end{array}$$

THEOREM 4. *If  $\alpha$  is of the form  $L^jM^kp \supset YL^{n+1}M^np$  ( $j \leq k$ ) then  $B + \alpha = T_n^+$  (and hence  $\alpha$  is a  $B^*$  formula).*

PROOF. Assume  $B$  and (1)  $L^jM^kp \supset YL^{n+1}M^np$ . From this and (1) we have  $p \supset YL^{n+1}M^np$ ; hence by Lemma 1,  $B + (1)$  contains  $T_n^+$ . T8 shows the converse.

THEOREM 5. *If  $\alpha$  is of the form  $XMp \supset YLp$  then  $B + \alpha = Triv$  (and hence, since  $Triv = T_0^+$ ,  $\alpha$  is a  $B^*$  formula).*

PROOF. Where  $X$  is any modality, the following are Theorems of  $B$ :

- $$\begin{array}{ll}
(1) & XM(p \supset Lp) \\
(2) & XL(p \supset Lp) \supset X(Mp \supset LLp)
\end{array}$$

Assume  $B$  and  $XMp \supset YLp$ . Then by  $[p \supset Lp/p]$  and (1) we have

- $$(3) \quad YL(p \supset Lp)$$

Let  $Y$  be of degree  $k$ . If  $k \leq 1$ , (3) yields  $p \supset Lp$  by T2 or T3. If  $k > 1$ , then

- $$\begin{array}{ll}
(4) & M^kL(p \supset Lp) \quad ((3), T5) \\
(5) & M^k(Mp \supset LLp) \quad ((4), (2)) \\
(6) & L^kMp \supset M^kLLp \quad ((5), T7) \\
(7) & M^kLL(p \supset Lp) \quad ((6)[p \supset Lp/p], (1)) \\
(8) & M^{k-1}L(p \supset Lp) \quad ((7), T2)
\end{array}$$

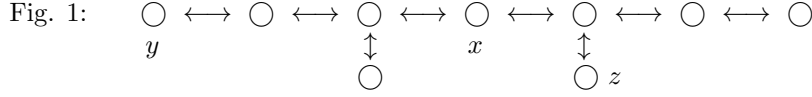
Whence by repetition we have  $ML(p \supset Lp)$ , and so by T2,  $p \supset Lp$ . Hence  $B + XMp \supset YLp$  contains  $Triv$ ; and T8 yields the converse.

It might be thought that all *wff* of the form  $Xp \supset Yp$  would turn out to be  $B^*$  formulae. This, however, is not the case.

THEOREM 6. *The logic  $B + (C) LMMp \supset MLLMMp$*

- $$\begin{array}{ll}
(a) & \text{is not contained in } B, (b) \text{ is contained in } T_2^+, \\
(b) & \text{does not contain any } T_n^+.
\end{array}$$

PROOF. (b) holds by T8. For (a): every reflexive symmetrical frame is a frame for  $B$ , but  $(C)$  is false at  $x$  in the model on the (reflexive) frame of Fig. 1 in which  $(V(p) = [y, z])$ .



For (c), define a *chain* as a linear reflexive symmetrical frame, i.e. as  $(W, R)$  where  $W = \{u_1, \dots, u_j, \dots\}$  and  $u_j R u_k$  iff  $k = j$  or  $k = j + 1$  or  $k = j - 1$ . It is not hard to show that every chain is a frame for  $B + (C)$ . But in an  $n + 2$ -membered chain  $L^n p \supset L^{n+1} p$  is false at  $u_1$  if  $p$  is false at  $u_{n+2}$  and nowhere else.

THEOREM 7. *The logic  $B + (D)LLLMMp \supset MMMLLMp$*

*(a) is contained in  $S5$ , (b) is independent of  $T_2^+$ ;  
and (c)  $T_2^+ + (D) = S5$ .*

PROOF. (a) holds by T8. For (c):  $MMp \supset LLLMMp$  and  $MMMLLMp \supset Mp$  are theorems of  $T_2^+$ . Hence  $T_2^+ + (D)$  yields  $MMp \supset Mp$ , and therefore  $S5$ ; and by (a) we have precisely  $S5$ . For (b): since  $T_2^+$  is weaker than  $S5$ , (c) shows that  $(D)$  is not in  $T_2^+$ . Moreover  $(D)$  has the striking property of being valid on a 4-membered chain, though not on one with 3 or more than 4 members. But as shown above,  $LLp \supset LLLp$  is invalid on a 4-membered chain, and so  $B + (D)$  does not contain  $T_2^+$ .

Note that  $(D)$ , unlike  $(C)$ , is valid on the frame of Fig. 1. Since  $(C)$ , unlike  $(D)$ , is valid on every chain,  $B + (C)$  and  $B + (D)$  are mutually independent.

Theorems 6 and 7 will generalize as follows: Let  $(C_m) = L^{m-1}M^m p \supset ML^m M^m p$ . Then if  $m \geq 2$ ,  $B + (C_m)$  is not contained in  $B$ , is contained in  $T_m^+$ , but does not contain any  $T_n^+$ . Let  $(D_m) = L^{m+2}M^{m+1}p \supset M^{m+2}L^{m+1}M^m p$ . Then if  $m \geq 1$ ,  $B + (D_m)$  is contained in  $T_m^+$  but is independent of  $T_{m+1}^+$ , and  $T_{m+1}^+ + (D_m) = T_m^+$ . (Note: each  $(D_m)$  is valid on an  $M + 3$ -membered chain.)

The following questions suggest themselves: (1) What classes of frames, if any, characterize the  $B + (C_m)$  and the  $B + (D_m)$  logics? (2) Does any

*wff* of the form  $Xp \supset Yp$ , when added to  $B$ , give either a logic between  $B + (D)$  and  $S5$  or one independent of both  $T_2^+$  and  $B + (D)$ ?

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