

M. W. Bunder

## A NOTE ON QUANTIFIED SIGNIFICANCE LOGICS

In 7.5 of [3] Goddard and Routly propose two logical systems of predicate calculus each incorporating certain significance restrictions.

In the first *PM1*, they read, in this connection ‘ $\vdash A$ ’ as ‘the significance restriction of  $A$  is a theorem’, where those variables in  $A$  that have free occurrences in  $A$  are restricted to their appropriate significance ranges. Using this interpretation of  $\vdash$  they then use the usual form of modus ponens:

$$\underline{R1} \quad \vdash A, \vdash A \supset B \rightarrow \vdash B,$$

which clearly fails to hold here.

Consider for example the case where  $A$  and  $B$  both contain the free variable  $u$  and that  $A$  is significant if  $u \in X_A$  (the significance range of  $A$ ) while  $B$  is significant if  $u \in X_B$ . Then  $A \supset B$  is significant if  $u \in X_A \cap X_B$  and *R1* will only give us  $\vdash B$  for  $u \in X_A \cap X_B$ .

Another interpretation of ‘ $\vdash A$ ’ namely: ‘if  $A$  is derivable independently of significance considerations, then  $A$  is a thesis for its significant values only’, also does not seem satisfactory as it leaves open the possibility of having a particular substitution instance of  $A$  derivable but not a thesis. Even if *SA* ( $A$  is significant) could be derived for this  $A$ , *SA* may also be a nonthesis.

In a second system *PM2*, they replace  $\vdash$  by  $\in$ , where ‘ $\in A$ ’ reads the universal closure of  $A$  is a theorem (where of course, each universal quantifier ranges of the significance range of the appropriate free variable).

Their modus ponens now becomes:

$$\underline{R1} \quad \in A, \in A \supset B \rightarrow \in B.$$

Taking the one variable case again,  $\in A$  means  $A$  holds for all  $u \in X_A$ ,  $\in A \supset B$  means that  $A \supset B$  holds for all  $u \in X_A \cap X_B$  so that we have  $B$  for  $u \in X_A \cap X_B$  rather than  $X_B$ .

If in this system, or even some part of it a given variable  $u$  has a range  $X_u$  over which all the predicate involving  $u$  in the (part) system are significant, this problem can be resolved. ' $\vdash A$ ' can then be read as ' $A$  with its free variables  $u_1, u_2, \dots, u_n$  restricted to  $X_{u_1}, X_{u_2}, \dots, X_{u_n}$  is a theorem'. Of course in general  $X_{u_1}$  will only be a subset of the significance range of  $u_1$  in any particular formula.

An alternative solution involves retaining  $R1$  only for cases where  $A$  and  $B$  have no free variables and using, for the one free variable case, the restricted generality of [2]. Writing  $A(u)$  and  $B(u)$  for  $A$  and  $B$  of  $R1$  we then have:

$$\vdash A(V), \vdash A(u) \supset_u B(u) \rightarrow \vdash B(V)$$

where ' $A(u) \supset_u B(u)$ ' is read ' $B(u)$  holds for all  $u$  for which  $A(u)$  holds' and where  $V$  is a term for which  $A(V)$  is derivable (as well as significant).

We could in fact assume:

$$D \vdash S(D)$$

i.e. whenever  $D$  is derivable,  $D$  is significant.

In the case of several free variables  $u_1, u_2, \dots, u_n$ , the generalized restricted generality of [1] can be used and we have instead of  $R1$ :

$$\begin{aligned} \vdash A(V_1, V_2, \dots, V_n), \vdash A(u_1, u_2, \dots, u_n) \supset u_1, u_2, \dots, u_n^{B(u_1, u_2, \dots, u_n)} \\ \rightarrow \vdash B(V_1, V_2, \dots, V_n). \end{aligned}$$

Of the quantifiers proposed in [3], those restricted to significance ranges can also be conveniently represented using restricted generality.

$(u)A(u)$ , where  $u$  is understood to range only over those values for which  $A(u)$  is significant, we can in fact define by  $S(A(u)) \supset_u A(u)$ .

Generalization ( $R2$  in [3]) then becomes:

$$S(A(u)) \vdash A(u) \rightarrow (u)A(u)$$

(where the  $\vdash$  is now interpreted in the usual fashion), and the remaining axioms and false of  $PM1$  could become:

*P1* If  $A$  is a substitution instance of an  $S_0$ -tautology then  $SA \vdash A$ .

*P2*  $SA, S((u)B) \vdash (u)(A \vee B) \supset .A \vee (u)B$ .

*P3*  $S(A(V)) \vdash (u)A(u) \supset A(V)$ .

Clearly the Behmann formulae such as

$$(x)f(x) ((x)f(x) \supset g(r)) \supset (x)g(x)$$

are not provable under this definition of  $(x)$ , which is as required.

## References

- [1] M. W. Bunder, *Generalized generality*, **Notre Dame Journal of Formal Logic**, Vol. XX (1979), pp. 620–624.
- [2] H. B. Curry and F. Feys, **Combinatory Logic**, North Holland, Amsterdam 1958.
- [3] L. Goddard and R. Routley, **The Logic of Significance and Context**, Scottish Academic Press, Edinburgh 1973.

*Department of Mathematics  
The University of Wollongong  
Wollongong, N. S. W., Australia*